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Synoptic Tables for the Solution of Ciphers

and

A Bibliography of Cipher Literature

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RIVERBANK LABORATORIES
 DEPARTMENT OF CIPHERS
 RIVERBANK
 GENEVA, ILL.
 1918

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by
William F. Friedman

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cations Security Section,
Office of Naval Communications.
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FOREWORD

The tables presented herewith are designed to meet specific pedagogical needs of a course of instruction in modern ciphers. They are not intended, it is frankly admitted, to serve as a guide for the expert in his attempt to analyze complex ciphers such as may be intercepted today.

The method which has been followed in their construction is analogous with that followed in chemical analysis manuals, but only in its broader aspects. The basis for the chemical determination of the nature of an unknown substance consists in the ability to place the unknown successively into one of two alternative classes by means of a series of definite tests until with the last cleavage the solution is reached. It is entirely possible to accomplish this determination with directness and with accuracy in chemical analysis because the laws underlying chemical reactions are definite and unchanging. The tests to be applied are exact, the reagents are all thoroughly understood. It is possible to determine the nature of even the most minute traces of an unknown substance, so refined have the methods of chemical analysis become. Contrast this situation with that which confronts the cipher analyst at the outset of his attempts to solve an unknown. In the first place, except in rare instances in practice, the amount of the unknown is often so limited as to thwart all his attempts at analysis and nothing can be done. In the second place, while it is true that both an unknown chemical substance and a message are composed of definite combinations of discrete units, the former of atoms, the latter of letters, further analogy between them ceases. For while atoms combine in a limited number of ways and positions to form molecules, and the latter combine in a limited number of ways and positions to form more complex substances, letters combine in a limitless number of ways and positions to form words, and words combine in a limitless number of ways and positions to form sentences. True, this difference is only one of degree, not of kind, but whereas the science of chemistry has reached so high a degree of development that each one of the possible combinations may be recognized by at least one test, the science and art of deciphering has not reached such a high level of perfection. In the field of complex ciphers, there is at present no definite means of determining what tests or what methods of solution should be applied because there is no way of determining from external characteristics or even from certain internal signs which one of a great number of complicated and readily modifiable systems of enciphering has been used in the particular message under examination. In fact, in most cases, unless the decipherer is able to secure some information concerning the system used he has no way of knowing what methods to apply until the long and laborious process of elimination has disclosed them.

The analogy between the tables for chemical and for cipher analysis is, therefore, only remote, and it is doubtful whether it can ever be brought closer. But for the purposes for which the tables presented are specifically intended, namely, instruction, it is believed that they will constitute a valuable adjunct to the curriculum of a course in ciphers. It is believed that they will afford the student a means of surveying the most important branches of practical ciphers and to note their similarities and differences. Thus, taken as a whole, they will give a more or less comprehensive bird's-eye view of the entire field of ciphers. If they will thus enable the student to secure a firmer grasp upon the basic principles underlying this branch of knowledge they will have served the purpose for which they were intended.

The Riverbank Publications referred to in the tables are as follows:

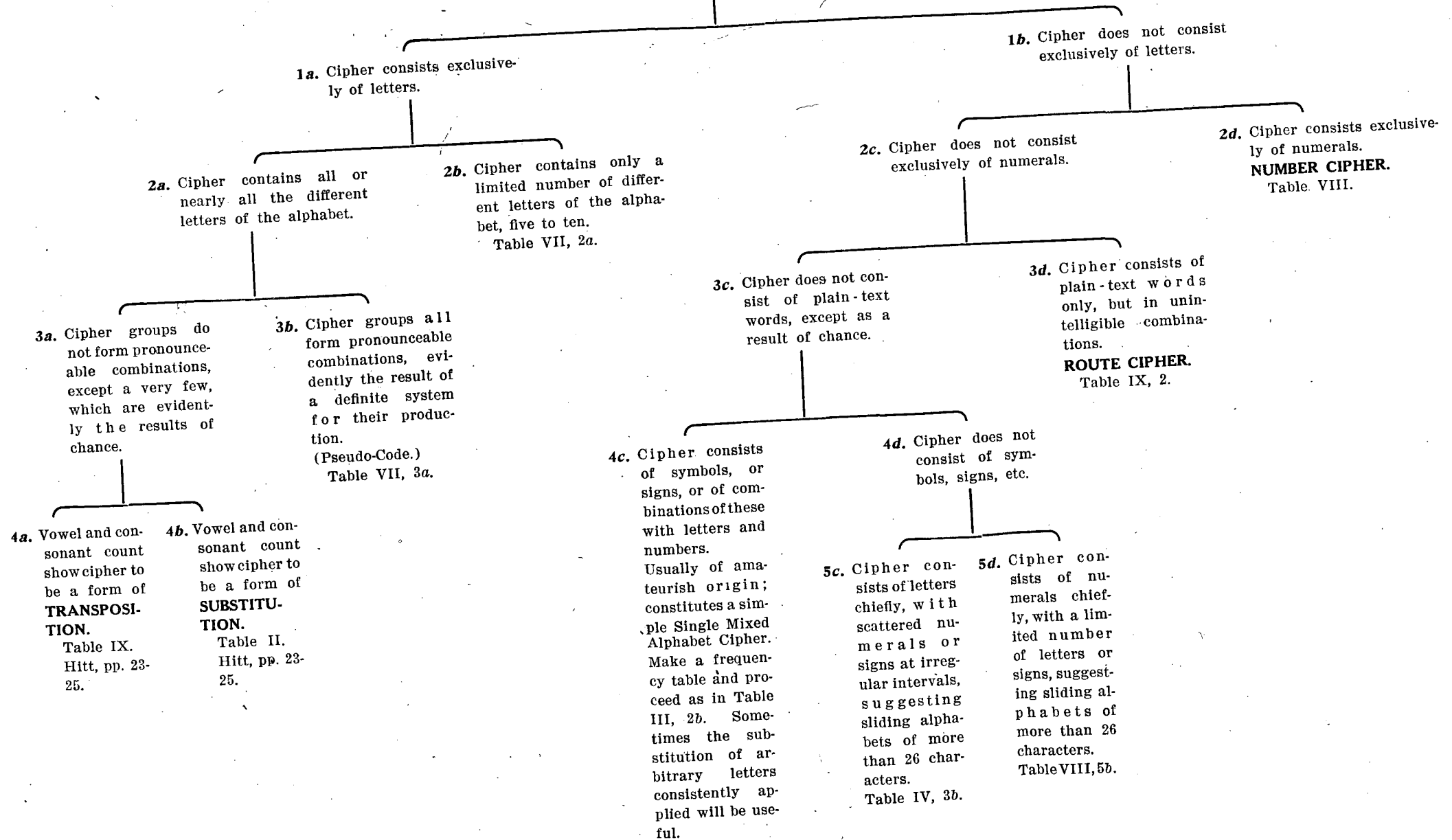
- No. 15—A Method of Reconstructing the Primary Alphabet given a Single One of the Series of Secondary Alphabets. 1917.
- No. 16—Methods for the Solution of Running-Key Ciphers. 1918.
- No. 17—An Introduction to Methods for the Solution of Ciphers. 1918.
- No. 19—Formulae for the Solution of Geometrical Transposition Ciphers.* 1918.
- No. 20—Several Machine Ciphers and Methods for their Solution.* 1918.
- No. 21—Methods for the Reconstruction of Primary Alphabets, Arbitrarily-Mixed Alphabets, Numerical Keys, etc.*

The full titles of works, which in the tables are referred to by only the author's name, will be found in the Bibliography, Part II, pages 14-16.

*Now in press.

TABLE I

Examine the cipher carefully in order to secure from extraneous circumstances such information as may be of value in the subsequent analysis. Certain clues may be found as to language, subject, correspondents, etc.



1. SUBSTITUTION CIPHER

Set a few groups on the Poly-Alphabet or apply the "running down" process.

[Also from Table VIII, 4a]

2a. Cipher solvable on the Poly-Alphabet, in the case of a single Straight Alphabet, or in the case of a series of Straight Alphabets wherein the words reappear on different lines.

Table III, 2a.

2b. Cipher not solvable on the Poly-Alphabet.

Apply the process of factoring the intervals separating recurring polygraphs, trigraphs, and digraphs.

3a. Factoring discloses no repeatedly recurring factors.

3b. Factoring discloses certain repeatedly recurring factors. (PERIODIC MULTIPLE ALPHABET SYSTEM.)

Table IV, 3a.

4a. Substitution equilateral, i.e., the total number of cipher letters is equal to the total number of plain-text letters.

4b. Substitution not equilateral, i.e., total number of cipher letters greater than total number plain-text letters, usually a multiple of the latter.

Table VII.

5a. Substitution monographic, i.e., letter for letter substitution, each one enciphered independently.

5b. Substitution not monographic.

6c. Substitution digraphic, i.e., pair for pair substitution.

6d. Substitution not digraphic.

6a. Frequency Table shows "crests and troughs."

SINGLE ALPHABET (MONO-ALPHABET) SYSTEM.

Table III.

6b. Frequency Table shows no marked "crests and troughs" but is "solid."

NON-PERIODIC MULTIPLE ALPHABET (POLY-ALPHABET) SYSTEM.

Table IV, 2b.

7a. PLAYFAIR SYSTEM

7b. Substitution by means of a rectangle.

Pages 5-8.

7c. Substitution Trigraphic. Pages 8-9.

7d. Substitution Polygraphic. (Approaching Code.)

8a. ORIGINAL PLAYFAIR SYSTEM

8b. MODIFIED PLAYFAIR SYSTEM

Solve by Mauborgne or Moorman method.*

Solve by combination of Mauborgne and Moorman method.*

* See Mauborgne, J. O., *An Advanced Problem in Cryptography and Its Solution*. Leavenworth Press, 1914. Hitt, *Manual for the Solution of Military Ciphers*, 1st edition, pp. 76-83; 2nd edition, pp. 76-82.

REF ID:A4146440

TABLE III

[From Table II, 2a and 6a; VII, 5a]

1. SINGLE ALPHABET (MONO-ALPHABET) SYSTEM

(Frequency table shows "crests and troughs")

[Also from Table I, 4c; VII, 2a; VIII, 5a; VIII, 5d]

2a. STRAIGHT ALPHABET CIPHER
(This should have been solved under Table II, 2a.)

2b. MIXED ALPHABET CIPHER

3a. DIRECT ALPHABET

3b. REVERSED ALPHABET

3c. RECIPROCAL ALPHABET

3d. NON-RECIPROCAL ALPHABET

4a. Solve by the Frequency Table Method, i.e., "fitting the frequency table to the normal" to find A. Substitute the plain-text values in sequence.

4b. Solve by means of a Poly-Alphabet or by applying the "running down" process.

4c. Solve by the Frequency Table Method, same as in 4a of this table.

4d. Find the Reversed Alphabet equivalents for three or four groups and proceed as in 4b of this table. Find the key letter and apply to the entire message.

See Riverbank Publication No. 17, pages 25-36, and Hitt, pages 39-62.

Make a frequency table with prefixes and suffixes and assume values based upon the frequency of individual letters, digraphs, and trigraphs.

See Riverbank Publication No. 17, pages 37-46, and Hitt, pages 44-50.

In case of a Reciprocal Alphabet, assignment of values is aided by the reciprocal relation. If the deciphering alphabet when completed exhibits signs of its being a Secondary Alphabet, based upon a Primary Alphabet using a key word, reconstruct the Primary Alphabet; or if the deciphering alphabet when completed exhibits signs of its being derived from a generating rectangle, reconstruct the latter. Sometimes these operations, when attempted upon the basis of partially deciphered material, will result in the complete reconstruction of the alphabet and the consequent entire decipherment. See Riverbank Publications Nos. 15, 16, and 21.

TABLE V

1. MULTIPLE ALPHABET SYSTEM—Continued

[From Table IV, 4a]

(Periodicity governed by the successive single letters of the plain text.)

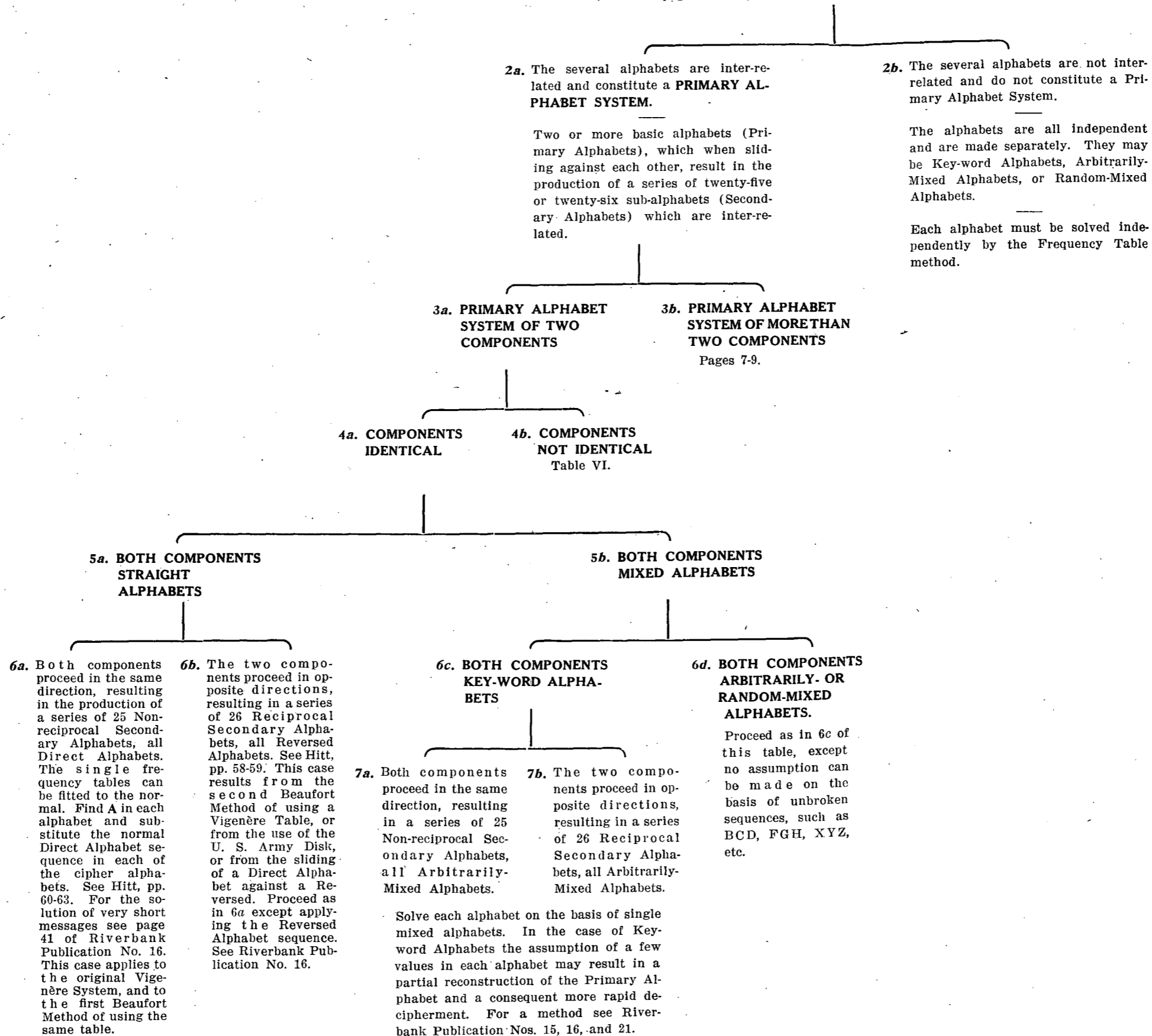


TABLE VI

[From Table V, 4b]

1. MULTIPLE ALPHABET SYSTEM—Continued

2. Primary Alphabet System of two components which are not identical.

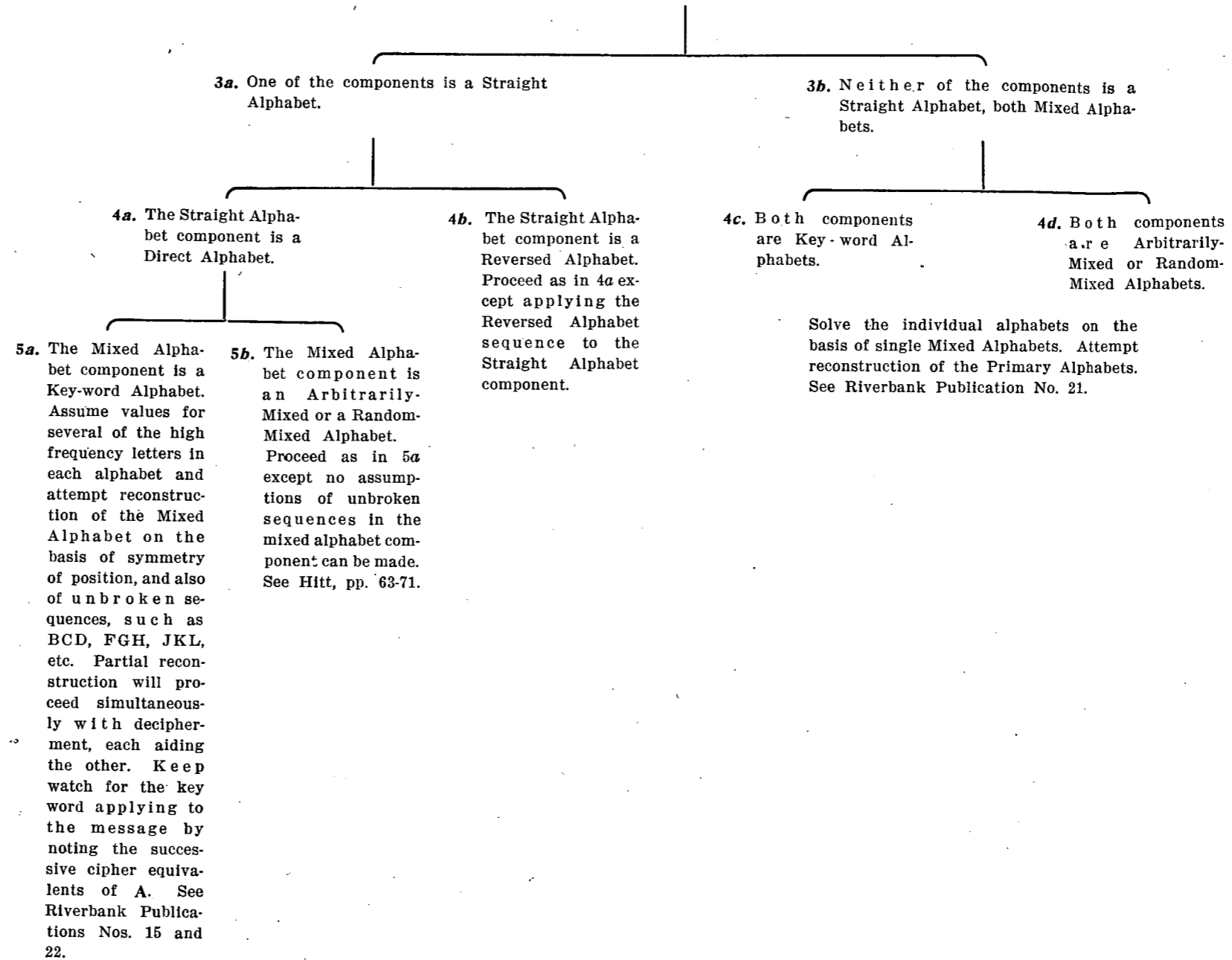


TABLE VII

[From Table II, 4b]

1. SUBSTITUTION NOT EQUILITERAL

Usually, if the number of plain-text letters is n , the number of cipher letters is $2n, 3n$, etc.

[Also from Table I, 2b]

2a. The number of different letters in the cipher message limited, usually not more than ten.

Systems using alphabets consisting of the various combinations of 2, 3, 4 . . . elements. (See Myers, pp. 65-165.) The number of characters in each combination is determined by the number of elements in the system. The least number of combinations possible must approximate 26, one for each letter of the alphabet.

- $2^n = 2^5 = 32$, a Biliteral Alphabet
- $3^n = 3^3 = 27$, a Triliteral Alphabet
- $4^n = 4^3 = 64$, a Tetraliteral Alphabet
- $5^n = 5^2 = 25$, a Pentaliteral Alphabet
- etc., etc.

Example of a Pentaliteral Alphabet, resulting from the use of a rectangle and a key word:

	G	R	A	N	T
G	A	B	C	D	E
R	F	G	H	I	K
A	L	M	N	O	P
N	Q	R	S	T	U
T	V	W	X	Y	Z

Example:

Plain text—T H E
Cipher—NW RA GT

Solution: Make a frequency table of combinations, or assign arbitrary single letters to each different combination and then make a frequency table. Proceed as in Table III, 2b. See Hitt, pages 83-85.

2b. The number of different letters in the cipher message approximates 26.

[Also from Table I, 3b]

3a. Cipher groups all pronounceable. (Pseudocode.)

4a. Regular arrangement of vowels and consonants, of the form CVCVC or VCVCV; groups all of the same length, either 5 or 10 letters.

4b. No regular arrangement or alternation of vowel and consonant.

Syllable ciphers (Built-up ciphers). Groups of irregular lengths usually. Substitution of letters or syllables for syllables of the plain text. Not often found and difficult to decipher in case of good systems. Solution by frequency of digraphs and trigraphs of the language.

3b. Cipher groups not pronounceable, except as the result of chance. Cipher groups usually the result of a square or a rectangular table, for enciphering not only letters but also syllables, words, phrases, etc. Approaching a code system.

The alphabets at the sides may be Key-Word Alphabets, Arbitrarily-Mixed, or Random-Mixed Alphabets, or numbers.

Solution: Make a frequency table of pairs and apply the frequency table method modified by considerations arising from the frequency of the most common words, for which substitution will have to be made by single pairs. Attempt a reconstruction of the alphabets at the sides.

5a. Regularity produced by the insertion of nulls. Compile a frequency table on the basis of every other letter and proceed as in Table III.

5b. Regularity produced by means of a table, or a rectangle on one side of which are consonants only, on the other side vowels only. Each plain-text letter requires two cipher letters.

Compile a frequency table of pairs and attempt reconstruction of the table, or rectangle.

6a. Rectangle based upon Straight Alphabets only. Make a frequency table of pairs and attempt a reconstruction of the rectangle.

6b. Rectangle not based upon Straight Alphabets.

7a. Rectangle based upon a Key-word Alphabet.

7b. Rectangle based upon an Arbitrarily-Mixed or a Random-Mixed Alphabet.

Solution by frequency of pairs. Attempt reconstruction of table.

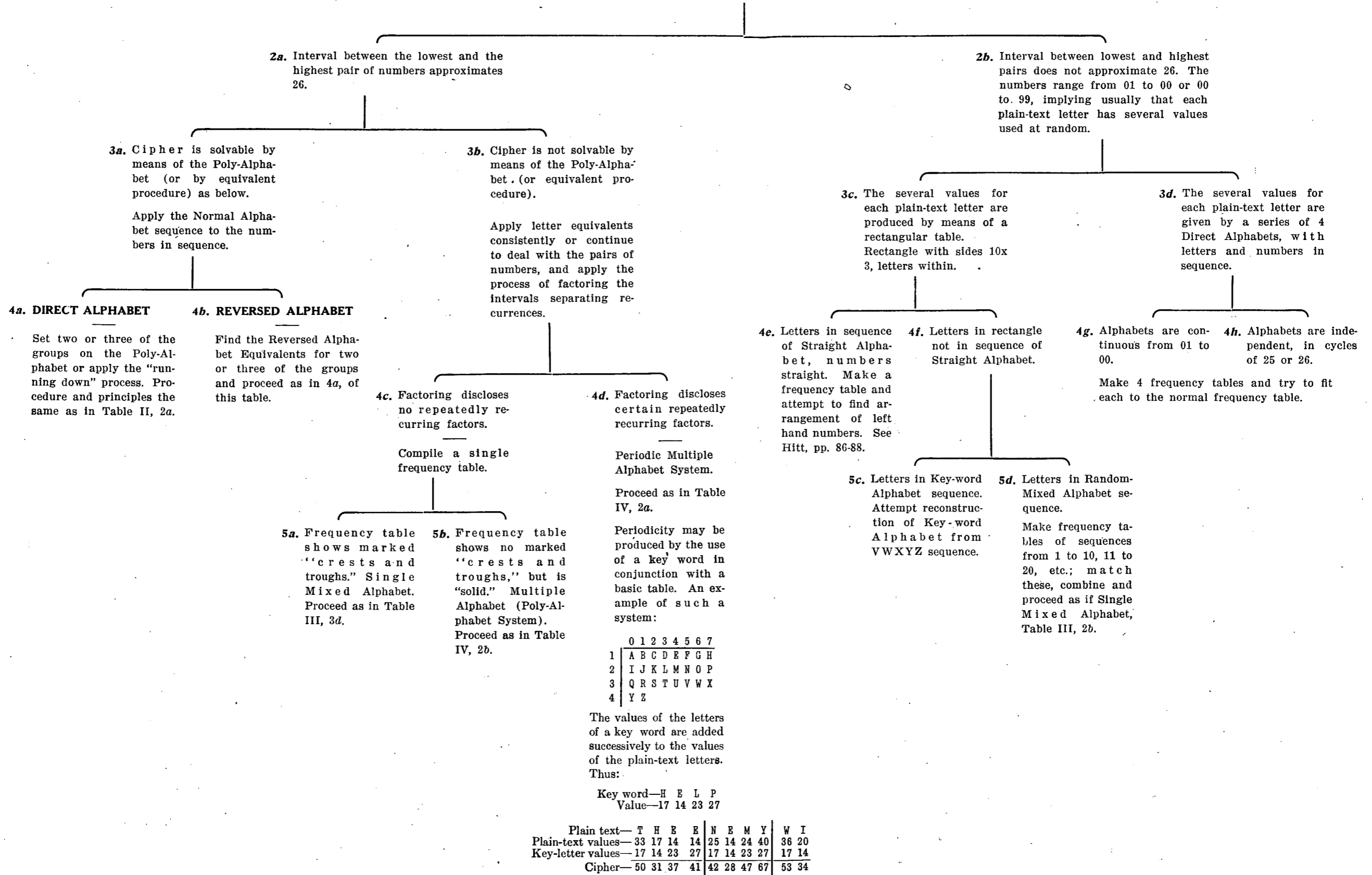
TABLE VIII

[From Table I, 2d]

1. NUMBER CIPHER

(Mathematical Ciphers)

Divide up the message into pairs of numbers unless already in this form.



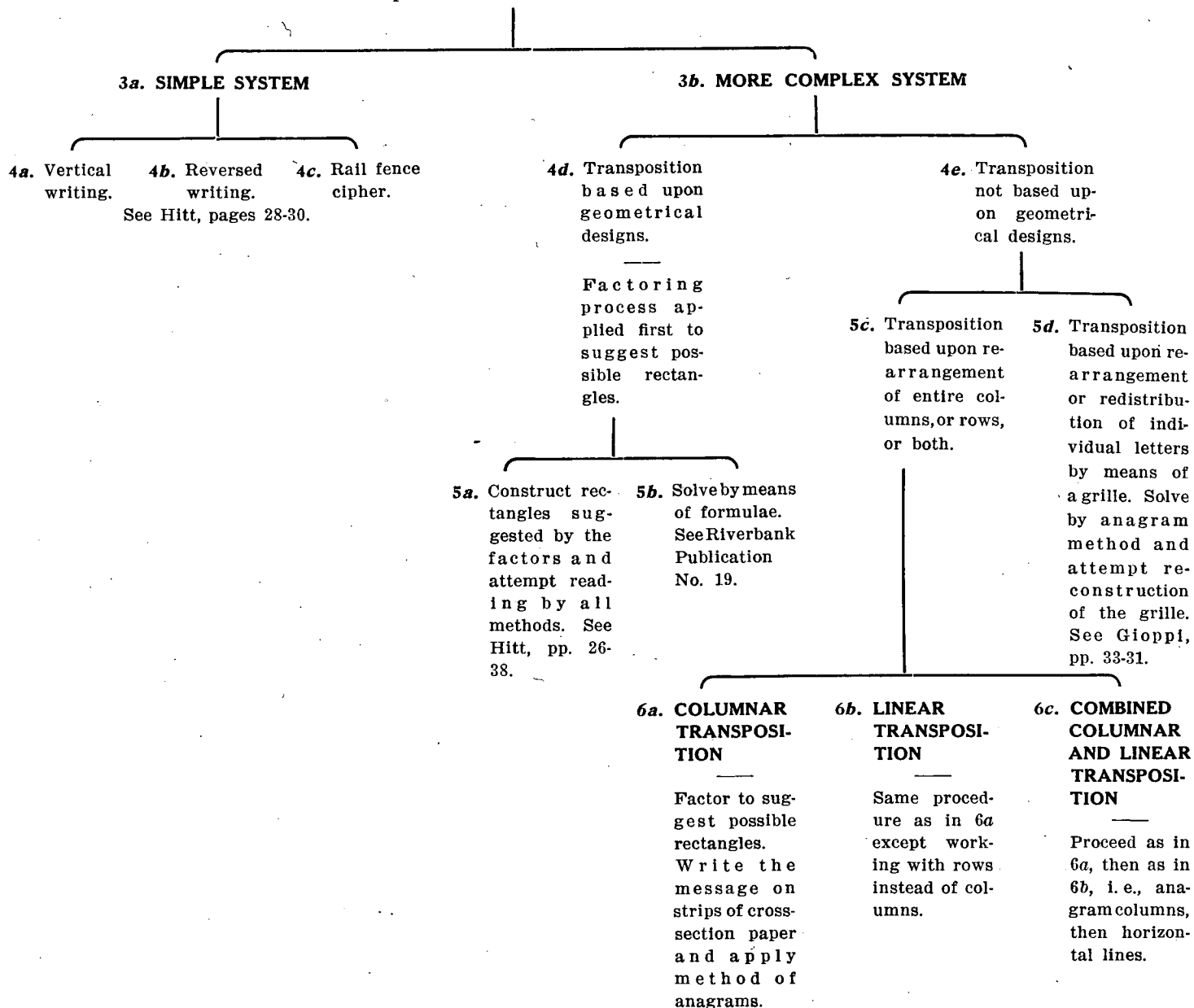
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TABLE IX

[From Table I, 4a]

1. TRANSPOSITION CIPHER

2. Including Route Ciphers, which are only a type of transposition ciphers wherein the words are treated as individual letters. Regard each word as a single letter or apply arbitrary letters or numbers to the words and proceed as below.



See Hitt, pp. 26-38.

DIGRAPHIC AND TRIGRAPHIC SUBSTITUTION

The chief advantage of digraphic and trigraphic substitution is that it prevents the decipherer from basing his analysis upon the frequency of individual letters in the language, and forces him to base any analysis to be made upon the frequency of digraphs and trigraphs: a circumstance which causes the analysis to become correspondingly difficult and, in addition, lessens the reliance which may be placed in it.

There are several ways of procuring digraphic substitution, of which the Playfair System is by far the most practical. Most of the other systems require tables, the use of which entails the expenditure of much labor, and the loss of one copy of which renders the entire system utterly unsafe. An excellent example of such a table is that shown in Fig. 1, which was taken from *La Crittografia*, pp. 84 and 85. Here the reciprocal relation

	+	A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	X	Y	Z	W	
+	ix	gu	do	mq	ag	jh	sp	cf	ki	bz	yq	ém	qr	kk	fp	kn	uo	ec	hj	ip	df	nc	mr	ct	di	md	ha	+
A	vb	xz	kj	yj	hp	plet	+d	ci	dw	xn	zly	pv	hh	cc	rf	ex	ps	zy	hy	sr	yo	fb	gn	wg	ij	A		
B	ek	lp	qt	hors	ur	cr	zh	gv	wchl	yn	kt	wt	me	kh	ev	gf	sq	yt	wa	hv	fn	vo	eu	+i	gd	B		
C	gi	dx	mu	ao	nh	sf	+g	wl	mm	ah	gr	b+	hs	zu	ym	wur	rz	ey	bf	ta	+x	ld	qb	aq	vt	qu	us	C
D	xk	km	yz	ry	f+	tr	+t	xo	jk	+y	po	gj	jx	pe	mo	+b	ak	hw	xur	hk	yi	nq	ca	lf	wy	ai	D	
E	lw	qu	hp	qg	jq	+q	ob	san	lpx	op	vs	af	+k	xr	u+	nt	tz	li	ra	kd	by	sl	zg	cq	ju	bp	E	
F	dd	mi	ax	nb	wj	lc	zs	io	uar	rv	s+	tx	oy	jc	bv	tt	+n	lo	tg	rq	vz	ls	gs	yy	jl	hn	nv	F
G	ke	yk	rm	vz	oa	ov	bq	yi	xac	+d	ky	nw	tq	ay	zm	rr	yb	ej	fv	on	+a	bh	tu	sz	me	kh	G	
H	qa	+w	k+	pt	to	bc	xq	lr	an	vq	+r	vp	bj	yx	fz	lz	eb	ad	wd	cl	qo	pn	bu	vw	at	qf	dq	H
I	mf	oh	tn	ux	ue	fg	qc	tb	om	du	aw	rt	xe	vy	qw	ya	+s	oz	qv	ug	pq	vh	tj	++	qz	xy	ou	I
J	yv	re	wk	fm	ty	zo	ka	o+	+o	kb	xs	dh	fy	ql	vf	uv	ok	edy	op	rv	go	qz	dl	mz	kl	uy	J	
K	hb	jj	ji	g+	et	su	xo	v1	bo	+h	ab	+m	jz	da	+o	sw	vm	sg	um	yc	bl	vt	xd	gw	dt	yh	qo	K
L	sm	uj	pw	fe	cu	wb	dy	uj	vk	ér	nf	wn	r1	sd	oe	fq	ban	ph	gf	sh	xg	uw	ms	pp	jh	oe	L	
M	rl	wv	ud	bn	+z	gz	i+	tw	wp	fa	nu	pu	pp	ch	qq	dn	vi	+c	+v	lx	of	cb	se	py	gk	ju	ru	M
N	te	pb	fc	+u	rg	xp	lj	so	ed	os	la	ut	eh	xw	pg	qi	lq	dv	+r	oe	pm	fw	uf	wo	xt	gl	N	
O	ig	gd	ef	vd	zn	ln	mt	rp	id	sn	wm	jp	rb	ih	gt	pe	ej	ju	uz	ni	vr	iw	ge	vv	fl	iq	zv	O
P	ws	ul	na	oo	sk	dm	yn	nn	q+	z+	ly	rf	ae	tv	hu	dj	ml	it	js	ar	hc	mk	xh	oi	mx	sj	lb	P
Q	ph	h+	cv	if	bw	kw	hz	ec	ze	no	vx	re	jm	tie	ah	tw	mn	+l	tm	bb	cz	ir	yd	uh	ty	in	Q	
R	uq	es	ol	ja	xi	qk	ap	nd	ds	ll	zk	zq	m+	gb	ys	ns	og	fs	gp	bd	ik	mw	fi	ve	dc	op	tf	R
S	fj	eg	zr	zw	lm	mv	ce	kq	lt	tk	pz	pd	ev	l+	oi	ng	+f	br	au	ub	zike	tc	yw	za	gy	ko	S	
T	nr	es	ig	sv	x+	n+	rw	fr	yr	qm	iv	si	wr	qs	ib	hd	vj	gm	dew	xf	og	pm	fk	jd	eq	mg	T	
U	eo	fh	ss	xl	mb	id	ux	is	qy	zj	lg	dp	pa	kr	wf	+b	zb	r+	be	cw	nk	zx	jo	ic	ju	or	lv	U
V	cy	ze	a+	xm	oc	yf	jn	jt	iu	mp	tp	lh	kg	kp	am	bx	bki	ot	ek	ku	y+	ox	qj	im	ft	hx	V	
X	td	gh	yg	dg	kv	il	wi	lu	pv	rd	zb	d+	uc	vc	aj	kf	ne	hf	en	jj	nz	dr	zz	wh	iz	aan	X	
Y	vu	io	gq	ks	qz	av	ve	xb	kz	gg	ac	ga	zd	cn	bj	ra	+j	th	rn	bs	pf	jt	km	fx	db	sx	Y	
Z	pi	sy	xj	qh	yl	va	pk	ex	bg	st	uir	jak	go	od	je	w+	rk	sb	ff	up	om	ow	uu	as	xv	sc	Z	
W	zp	bt	le	bi	hr	rx	un	az	xx	if	fd	jb	og	oj	lk	ny	mh	wq	tl	p+	bm	co	ma	ts	dz	ge	qp	W
+	A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	X	Y	Z	W		

Fig. 1

of plain text and cipher text is such that the same table can be used for enciphering and deciphering. For example:

Enciphering—TH EN EM YP RE PA RE
YR XR +K AL QK UL QK

Deciphering—YR XR +K AL QK UL QK
TH EN EM YP RE PA RE

Note that two pairs, even if they involve a common letter, do not have a common letter in the cipher equivalent, except as a matter of chance. The result of this fact is that no grouping of cipher pairs representing combinations of E with other letters can be made upon the basis of a common letter in such cipher pairs.

The process of arranging such a table, however, is very laborious, so that frequent change is impractical. Another form of such a table which may, on the other hand, be changed very frequently, but which does not possess the reciprocal relation, is that shown in Fig. 2, but here there is an added disadvantage—that of having a common cipher letter as a result in those pairs which represent plain-text pairs having a letter in common. Thus ER, EN, ES, and ET are enciphered by TU, TK, TV, and WT respectively, or by the reversals of the latter. These digraphs are found at the intersection of the vertical column determined by the first letter of each pair as located in the top line, and the row determined by the second letter of each pair as located in the first column at the left. When the cipher pair is taken at the intersection of the row determined by the first letter, and the vertical column determined by the second letter of each pair, the equivalents for these same combinations are UK, KF, VL, and WN, or their reversals; but note that all the combinations ending with the same letter will show a letter in common.

The same results may be obtained by employing sliding strips, as shown in the accompanying diagram. The direct alphabet, I, and the second mixed alphabet, IV, are fixed; the first mixed alphabet, III, is mounted upon a movable strip with another direct alphabet, II; the sliding alphabets are moved so that the first letter of the pair on alphabet II is placed beneath A on alphabet I, then under the second letter of the pair on I, the two cipher equivalents of the pair are found on III and IV. Thus, for the word THIS the successive positions and encipherments are as follows:

TH = S A	{	I—ABCDEFGHIJKLMN OPQRSTUVWXYZ	Fixed Alphabet
		II—TUVWXYZABCDEFGHIJKLMN OPQRS	} Movable Alphabets
		III—MQUVWXYZSTENOGRAPHYBCDFIJKL	
		IV—CRYPTOGAMSBDEFHIJKLNQUVWXZ	Fixed Alphabet
I S = S L	{	I—ABCDEFGHIJKLMN OPQRSTUVWXYZ	Fixed Alphabet
		II—IJKLMNOPQRSTUVWXYZABCDEFGHI	} Movable Alphabets
		III—PHYBCDFIJKLMQUVWXYZSTENOGRA	
		IV—CRYPTOGAMSBDEFHIJKLNQUVWXZ	Fixed Alphabet

	A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z
A	S	T	E	N	O	G	R	A	P	H	Y	B	C	D	F	I	J	K	L	M	Q	U	V	W	X	Z
B	C	R	Y	P	T	O	G	A	M	S	B	D	E	F	H	I	J	K	L	M	Q	U	V	W	X	Z
C	T	E	N	O	G	R	A	P	H	Y	B	C	D	F	I	J	K	L	M	Q	U	V	W	X	Z	
D	C	R	Y	P	T	O	G	A	M	S	B	D	E	F	H	I	J	K	L	M	Q	U	V	W	X	Z
E	N	O	G	R	A	P	H	Y	B	C	D	F	I	J	K	L	M	Q	U	V	W	X	Z	S	T	E
F	C	R	Y	P	T	O	G	A	M	S	B	D	E	F	H	I	J	K	L	M	Q	U	V	W	X	Z
G	O	G	R	A	P	H	Y	B	C	D	F	I	J	K	L	M	Q	U	V	W	X	Z	S	T	E	N
H	C	R	Y	P	T	O	G	A	M	S	B	D	E	F	H	I	J	K	L	M	Q	U	V	W	X	Z
I	G	R	A	P	H	Y	B	C	D	F	I	J	K	L	M	Q	U	V	W	X	Z	S	T	E	N	O
J	C	R	Y	P	T	O	G	A	M	S	B	D	E	F	H	I	J	K	L	M	Q	U	V	W	X	Z
K	H	Y	B	C	D	F	I	J	K	L	M	Q	U	V	W	X	Z	S	T	E	N	O	G	R	A	P
L	C	R	Y	P	T	O	G	A	M	S	B	D	E	F	H	I	J	K	L	M	Q	U	V	W	X	Z
M	B	C	D	F	I	J	K	L	M	Q	U	V	W	X	Z	S	T	E	N	O	G	R	A	P	H	Y
N	C	R	Y	P	T	O	G	A	M	S	B	D	E	F	H	I	J	K	L	M	Q	U	V	W	X	Z
O	D	F	I	J	K	L	M	Q	U	V	W	X	Z	S	T	E	N	O	G	R	A	P	H	Y	B	C
P	C	R	Y	P	T	O	G	A	M	S	B	D	E	F	H	I	J	K	L	M	Q	U	V	W	X	Z
Q	F	I	J	K	L	M	Q	U	V	W	X	Z	S	T	E	N	O	G	R	A	P	H	Y	B	C	D
R	C	R	Y	P	T	O	G	A	M	S	B	D	E	F	H	I	J	K	L	M	Q	U	V	W	X	Z
S	J	K	L	M	Q	U	V	W	X	Z	S	T	E	N	O	G	R	A	P	H	Y	B	C	D	F	I
T	C	R	Y	P	T	O	G	A	M	S	B	D	E	F	H	I	J	K	L	M	Q	U	V	W	X	Z
U	M	Q	U	V	W	X	Z	S	T	E	N	O	G	R	A	P	H	Y	B	C	D	F	I	J	K	L
V	C	R	Y	P	T	O	G	A	M	S	B	D	E	F	H	I	J	K	L	M	Q	U	V	W	X	Z
W	U	V	W	X	Z	S	T	E	N	O	G	R	A	P	H	Y	B	C	D	F	I	J	K	L	M	Q
X	C	R	Y	P	T	O	G	A	M	S	B	D	E	F	H	I	J	K	L	M	Q	U	V	W	X	Z
Y	W	X	Z	S	T	E	N	O	G	R	A	P	H	Y	B	C	D	F	I	J	K	L	M	Q	U	V
Z	C	R	Y	P	T	O	G	A	M	S	B	D	E	F	H	I	J	K	L	M	Q	U	V	W	X	Z

Fig. 2

Given a single long message or a series of messages in the same alphabets, a frequency table of pairs may be made the basis of solution, by assigning high-frequency-digraph values to the most frequent pairs. In the latter case, where two pairs having a common cipher letter have a common letter in their respective cipher equivalents, this relation would be a great aid in the assignment of values, since it would enable the decipherer to assign his values accordingly. In the case of key-word and direct alphabets the reconstruction of the alphabets may be attempted. Arbitrarily-mixed and random-mixed alphabets may also be used in such tables.

Still another form of table which may be used for digraphic substitution is that shown in Fig. 3. Here there are concerned one mixed and two direct alphabets and a

I—	ABCDEFGHIJKLMN	OPQRSTUVWXYZ
II—	FSKZRB	JEYQAHLTGXPDCUIWNVOM
III		
A	HTCGWSR	KBFJVIQAELUDPXMZOYN
B	TCGWSR	KBFJVIQAELUDPXMZOYNH
C	CGWSR	KBFJVIQAELUDPXMZOYNHT
D	GWSR	KBFJVIQAELUDPXMZOYNHTC
E	WSR	KBFJVIQAELUDPXMZOYNHTCG
F	SR	KBFJVIQAELUDPXMZOYNHTCGW
G	R	KBFJVIQAELUDPXMZOYNHTCGWS
H	K	BFFJVIQAELUDPXMZOYNHTCGWSR
I	B	FJVIQAELUDPXMZOYNHTCGWSRK
J	F	JVIQAELUDPXMZOYNHTCGWSRKB
K	J	VIQAELUDPXMZOYNHTCGWSRKB
L	V	IQAELUDPXMZOYNHTCGWSRKB
M	I	QAELUDPXMZOYNHTCGWSRKB
N	Q	AELUDPXMZOYNHTCGWSRKB
O	A	ELUDPXMZOYNHTCGWSRKB
P	E	LUDPXMZOYNHTCGWSRKB
Q	L	UDPXMZOYNHTCGWSRKB
R	U	DPXMZOYNHTCGWSRKB
S	P	XMZOYNHTCGWSRKB
T	X	MZOYNHTCGWSRKB
U	M	ZOYNHTCGWSRKB
V	Z	OYNHTCGWSRKB
W	O	YNHTCGWSRKB
X	Y	NHTCGWSRKB
Y	N	HTCGWSRKB
Z	H	TGWSRKB

Fig. 3

quadricular table. The first letter of a pair is sought in Alphabet I, its equivalent taken in Alphabet II, and by following the horizontal line in the quadricular table determined by the second letter of the pair in Alphabet III to the vertical column determined by the first letter, the cipher letter is taken at the intersection. Thus:

TH ER EI SN OT HI NG
UH RM RI CS GK EE TP

Note that as far as the first letter in each pair is concerned, the encipherment is merely by means of a single mixed alphabet. It is only the encipherment of the second letter which is multi-alphabetical in nature.

The same table shown in Fig. 3, with one additional alphabet, IV, may be used for trigraphic substitution. The equivalent of the first letter in a group is found in Alphabet II directly beneath that letter in Alphabet I. The equivalent of the second letter is found in Alphabet IV directly opposite the letter in Alphabet III. The equivalent of the third letter is found at the intersection of the horizontal line in the quadricular table determined by the second letter, and the vertical column determined by the position of the third letter in Alphabet I. Thus:

THE REI SNO THI NGT
URV DDI CQH URE TAN

The variations of this system are many; but as far as the two letters in each group of triplets is concerned, encipherment is purely mono-alphabetical. (See Gioppi, pp. 45-46.)

		I—ABCDEFGHIJKLMN	OP	QRST	UV	WXYZ							
		II—FSKZRBJEYQAHLTGXPD											
		CUIWNVOM											
III	IV												
A	K	HT	CG	WS	RK	BF	JVI	QA	EL	UD	PX	MZ	OYN
B	S	TC	GS	WR	KBF	JVI	QA	EL	UD	PX	MZ	OYN	H
C	B	CG	WS	RK	BF	JVI	QA	EL	UD	PX	MZ	OYN	HT
D	U	GS	WR	KBF	JVI	QA	EL	UD	PX	MZ	OYN	HT	C
E	D	WS	RK	BF	JVI	QA	EL	UD	PX	MZ	OYN	HT	CG
F	J	SR	KBF	JVI	QA	EL	UD	PX	MZ	OYN	HT	CG	W
G	A	RK	BF	JVI	QA	EL	UD	PX	MZ	OYN	HT	CG	WS
H	R	KBF	JVI	QA	EL	UD	PX	MZ	OYN	HT	CG	WS	R
I	V	BF	JVI	QA	EL	UD	PX	MZ	OYN	HT	CG	WS	RK
J	I	FJ	VI	QA	EL	UD	PX	MZ	OYN	HT	CG	WS	RKB
K	H	JVI	QA	EL	UD	PX	MZ	OYN	HT	CG	WS	RKB	F
L	T	VI	QA	EL	UD	PX	MZ	OYN	HT	CG	WS	RKB	FJ
M	L	IQA	EL	UD	PX	MZ	OYN	HT	CG	WS	RKB	FJ	V
N	Q	QA	EL	UD	PX	MZ	OYN	HT	CG	WS	RKB	FJ	VI
O	G	AE	LU	DP	XM	ZO	YN	HT	CG	WS	RKB	FJ	VI
P	C	EL	UD	PX	MZ	OYN	HT	CG	WS	RKB	FJ	VI	QA
Q	M	LU	DP	XM	ZO	YN	HT	CG	WS	RKB	FJ	VI	QA
R	F	UD	PX	MZ	OYN	HT	CG	WS	RKB	FJ	VI	QA	EL
S	X	DP	XM	ZO	YN	HT	CG	WS	RKB	FJ	VI	QA	EL
T	O	PX	MZ	OYN	HT	CG	WS	RKB	FJ	VI	QA	EL	UD
U	Y	XM	ZO	YN	HT	CG	WS	RKB	FJ	VI	QA	EL	UD
V	N	MZ	OYN	HT	CG	WS	RKB	FJ	VI	QA	EL	UD	PX
W	Z	ZO	YN	HT	CG	WS	RKB	FJ	VI	QA	EL	UD	PX
X	W	OYN	HT	CG	WS	RKB	FJ	VI	QA	EL	UD	PX	MZ
Y	P	YN	HT	CG	WS	RKB	FJ	VI	QA	EL	UD	PX	MZ
Z	E	NHT	CG	WS	RKB	FJ	VI	QA	EL	UD	PX	MZ	OY

Fig. 4

COMPLEX SYSTEMS

When the steps in analysis given in the preceding tables have failed to lead to results, it may be concluded that the cipher is either the result of (1) a modification or a combination of the systems enumerated, such as the combination of Substitution and Transposition systems, or (2) a system simple in itself as regards enciphering, but difficult in its results as far as deciphering is concerned. Some of the latter have been devised by experts who are in possession of all the known methods of attacking ciphers and have elaborated systems which allow no opening for the would-be decipherer. No attempt is made here to enumerate or to elucidate all of these systems, but among them may be mentioned the following:

- (1) Running Key Systems
- (2) Multiplex Alphabet Systems
- (3) Wheatstone Principle Systems
- (4) Fractionating Systems
- (5) Auto-key Systems
- (6) Variable Key Systems

(1) Running Key Systems. These systems make use of the running text of a book, identical copies of which are in possession of the correspondents. For a brochure on the subject see Riverbank Publication No. 16.

(2) Multiplex Alphabet Systems. These systems make use of a machine on the principle of the Bazeries disk cipher (Bazeries, pp. 250-261). For a brochure on the subject see Riverbank Publication No. 20; also De Viaris, "*L'Art de Chiffrer*," pp. 99-109.

(3) Wheatstone Principle Systems, which are based upon a mechanical cryptograph invented by Sir Charles Wheatstone in 1879. For a discussion of such a cipher and methods for solving it see Riverbank Publication No. 20.

(4) Fractionating Systems. The basic principle here is that the cipher letters or cipher numbers are compounded from parts of plain-text letters according to some definite system. A simple example is the following:

Alphabet—	A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z
Numerical Value—	01	02	03	04	05	06	07	08	09	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25	26

Each letter is represented by two digits. Write the dispatch horizontally, then apply the two digits for each letter one under the other. Thus:

ENEMY PREPARES
 01012 11010101
 54535 68561859

The cipher then is taken in any way in which a rearrangement of the digits may be effected. Thus, a very simple way would be to take the cipher digits in pairs from horizontal lines, and then find their letter equivalents on the conventional alphabet. This dispatch would begin

AAVJJ OSI etc.

In the case of any cipher number above 26, deduct 26 or a multiple thereof and find the equivalent of the remainder. Variations of the system are legion in number. The plain text may be written in groups of three, four, or five letters and the cipher letters may be selected accordingly upon some different scheme. This system, because of the number of unknown factors which are presented to the would-be decipherer, is a very difficult one to solve. Fractionating systems in which each cipher letter represents the halves, thirds, quarters, and possibly greater fractions of 2, 3, 4, or 5 plain-text letters may be devised, and would tax the ingenuity of the expert decipherer. (See Gioppi, pp. 102-114.)

(5) Auto-key Systems. Sometimes called Auto-enciphering Systems. This system was described by Vigenère, reinvented in 1884 by Captain Delauney, and perfected by Josse. The basic principle is that each cipher letter automatically becomes the key for the encipherment of the succeeding plain-text letter. Usually a key-word alphabet or a random-mixed alphabet is used, the letters of which are numbered in sequence. Thus:

AIWGHV L J X O C M Z P B K Y R D N T E Q U F S
 1 2 3 4 5 6 7 8 9 10 11 12 13 14 15 16 17 18 19 20 21 22 23 24 25 26

MESSAGE: Enemy prepares, etc.

	E	N	E	M	Y	P	R	E	P	A	R	E	S
	22	20	22	12	17	14	18	22	14	11	18	22	26
	22	16	12	24	15	3	21	17	5	6	24	20	20
CIPHER:	E	K	M	U	B	W	T	Y	H	V	U	N	N

Each cipher letter is produced in turn by finding the letter-value of the sum of the numerical equivalent of the preceding cipher letter and that of the plain-text letter to be enciphered; when this total exceeds 26, the latter amount is deducted and the letter-value of the remainder is taken as the cipher equivalent.

The great disadvantage of this system is that an error in one place produces errors in all the succeeding letters so that the recipient is caused to lose much time in the translation of a message which has many errors. A method which dispenses with the numerals is to construct a quadricular table from the alphabet as shown in Fig. 6.

	A	I	W	G	H	V	L	J	X	O	C	M	Z	P	B	K	Y	R	D	N	T	E	Q	U	F	S
A	I	W	G	H	V	L	J	X	O	C	M	Z	P	B	K	Y	R	D	N	T	E	Q	U	F	S	A
I	W	G	H	V	L	J	X	O	C	M	Z	P	B	K	Y	R	D	N	T	E	Q	U	F	S	A	I
W	G	H	V	L	J	X	O	C	M	Z	P	B	K	Y	R	D	N	T	E	Q	U	F	S	A	I	W
G	H	V	L	J	X	O	C	M	Z	P	B	K	Y	R	D	N	T	E	Q	U	F	S	A	I	W	G
H	V	L	J	X	O	C	M	Z	P	B	K	Y	R	D	N	T	E	Q	U	F	S	A	I	W	G	H
V	L	J	X	O	C	M	Z	P	B	K	Y	R	D	N	T	E	Q	U	F	S	A	I	W	G	H	V
L	J	X	O	C	M	Z	P	B	K	Y	R	D	N	T	E	Q	U	F	S	A	I	W	G	H	V	L
J	X	O	C	M	Z	P	B	K	Y	R	D	N	T	E	Q	U	F	S	A	I	W	G	H	V	L	J
X	O	C	M	Z	P	B	K	Y	R	D	N	T	E	Q	U	F	S	A	I	W	G	H	B	L	J	X
O	C	M	Z	P	B	K	Y	R	D	N	T	E	Q	U	F	S	A	I	W	G	H	V	L	J	X	O
C	M	Z	P	B	K	Y	R	D	N	T	E	Q	U	F	S	A	I	W	G	H	V	L	J	X	O	C
M	Z	P	B	K	Y	R	D	N	T	E	Q	U	F	S	A	I	W	G	H	V	L	J	X	O	C	M
Z	P	B	K	Y	R	D	N	T	E	Q	U	F	S	A	I	W	G	H	V	L	J	X	O	C	M	Z
P	B	K	Y	R	D	N	T	E	Q	U	F	S	A	I	W	G	H	V	L	J	X	O	C	M	Z	P
B	K	Y	R	D	N	T	E	Q	U	F	S	A	I	W	G	H	V	L	J	X	O	C	M	Z	P	B
K	Y	R	D	N	T	E	Q	U	F	S	A	I	W	G	H	V	L	J	X	O	C	M	Z	P	B	K
Y	R	D	N	T	E	Q	U	F	S	A	I	W	G	H	V	L	J	X	O	C	M	Z	P	B	K	Y
R	D	N	T	E	Q	U	F	S	A	I	W	G	H	V	L	J	X	O	C	M	Z	P	B	K	Y	R
D	N	T	E	Q	U	F	S	A	I	W	G	H	V	L	J	X	O	C	M	Z	P	B	K	Y	R	D
N	T	E	Q	U	F	S	A	I	W	G	H	V	L	J	X	O	C	M	Z	P	B	K	Y	R	D	N
T	E	Q	U	F	S	A	I	W	G	H	V	L	J	X	O	C	M	Z	P	B	K	Y	R	D	N	T
E	Q	U	F	S	A	I	W	G	H	V	L	J	X	O	C	M	Z	P	B	K	Y	R	D	N	T	E
Q	U	F	S	A	I	W	G	H	V	L	J	X	O	C	M	Z	P	B	K	Y	R	D	N	T	E	Q
U	F	S	A	I	W	G	H	V	L	J	X	O	C	M	Z	P	B	K	Y	R	D	N	T	E	Q	U
F	S	A	I	W	G	H	V	L	J	X	O	C	M	Z	P	B	K	Y	R	D	N	T	E	Q	U	F
S	A	I	W	G	H	V	L	J	X	O	C	M	Z	P	B	K	Y	R	D	N	T	E	Q	U	F	S

Fig. 6

Proceeding down the column determined by E (the first letter of the message) in the first horizontal line, to the line determined by the next plain-text letter N, the letter K, at the intersection, is taken as the cipher letter. Proceeding down the column determined by K in the first horizontal line to the line determined by E, the third plain-text letter, the cipher letter M, at the intersection, is taken as the cipher letter, etc. (See Gioppi, pp. 42-44.)

A method which is the equivalent to the quadricular table in its final results and which is easier to operate, makes use of two sliding strips bearing the alphabets; by shifting the lower strip so that the letter which becomes the key letter for the next encipherment, is placed beneath the letter immediately preceding the first letter in the alphabet concerned, the

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